

Find shortest path from *q* to *s* in *P* 



Find shortest path from *q* to *s* in *P* 



Find shortest path from *q* to *s* in *P s* may move



Find shortest path from *q* to *S* in *P* sites in *S* may move



Find shortest path from *q* to *S* in *P* sites in *S* may move



Find shortest path from *q* to *S* in *P* sites in *S* may move



Find shortest path from *q* to *S* in *P* 

![](_page_7_Figure_3.jpeg)

Find shortest path from *q* to *S* in *P* 

![](_page_8_Figure_3.jpeg)

Find shortest path from *q* to *S* in *P* 

![](_page_9_Figure_3.jpeg)

Find shortest path from *q* to *S* in *P* 

![](_page_10_Figure_3.jpeg)

Find shortest path from *q* to *S* in *P* 

![](_page_11_Figure_3.jpeg)

Find shortest path from q to S in P

- 1. When does VD(S) change?
- 2. How many changes?
- 3. Can we maintain VD(S) efficiently?

![](_page_12_Figure_6.jpeg)

Find shortest path from q to S in PMaintain MST(S), NN(S), etc.

- 1. When does VD(S) change?
- 2. How many changes?
- 3. Can we maintain VD(S) efficiently?

![](_page_13_Picture_6.jpeg)

VD(S): (geodesic) Voronoi diagram

- 1. When does VD(S) change?
- 2. How many changes?
- 3. Can we maintain VD(S) efficiently?

VD(S) has size  $\Theta(n + m)$ 

![](_page_14_Figure_6.jpeg)

VD(S): (geodesic) Voronoi diagram

- 1. When does VD(S) change?
- 2. How many changes?
- 3. Can we maintain VD(S) efficiently?

VD(S) has size  $\Theta(n + m)$ 

![](_page_15_Picture_6.jpeg)

- 1. When does VD(S) change?
- 2. How many changes?
- 3. Can we maintain VD(S) efficiently?

![](_page_16_Picture_5.jpeg)

- 1. When does VD(S) change?
- 2. How many changes?
- 3. Can we maintain VD(S) efficiently?

![](_page_17_Picture_5.jpeg)

![](_page_18_Figure_0.jpeg)

![](_page_19_Figure_0.jpeg)

Bisector  $B_{pq}$ : size:  $\Theta(m)$ #changes:  $\Theta(m^3)$ 

![](_page_20_Figure_2.jpeg)

Bisector  $B_{pq}$ : size:  $\Theta(m)$ #changes:  $\Theta(m^3)$ 

KDS:

O(m) size  $O(m^3)$  events

O(log m) time per event

![](_page_21_Picture_5.jpeg)

Bisector  $B_{pq}$ : size:  $\Theta(m)$ #changes:  $\Theta(m^3)$ 

KDS: *O*(*m*) size *O*(*m*<sup>3</sup>) events *O*(log *m*) time per event *O*(log *m*) time update

 $B_{pq}$ Þ۶

![](_page_23_Picture_0.jpeg)

![](_page_24_Picture_0.jpeg)

Voronoi center  $c_{pqs}$ : size:  $\Theta(1)$ #changes:  $\Theta(m^3)$ 

![](_page_25_Figure_2.jpeg)

Voronoi center  $c_{pqs}$ : size:  $\Theta(1)$ #changes:  $\Theta(m^3)$ 

KDS:

O(m) size  $O(m^3)$  events  $O(\log^2 m)$  time per event  $O(\log^2 m)$  time update

![](_page_26_Picture_4.jpeg)

VD (S): size:  $\Theta(n + m)$ #changes:  $\tilde{O}(n^3m^3)$  $\Omega(nm^3 + n^2m)$ 

![](_page_27_Figure_2.jpeg)

VD (S): size:  $\Theta(n + m)$ #changes:  $\tilde{O}(n^3m^3)$  $\Omega(nm^3 + n^2m)$ 

In  $\mathbb{R}^2$ : #changes:  $\tilde{O}(n^3)$ 

![](_page_28_Figure_3.jpeg)

VD (S): size:  $\Theta(n + m)$ #changes:  $\tilde{O}(n^3 m^3)$   $\Omega(nm^3 + n^2m)$ KDS: O(n + m) size  $\tilde{O}(n^3 m^3)$  events  $O(k \log^2 m)$  time per event

k = #affected neighbours

![](_page_29_Figure_3.jpeg)

### Overview

1.  $\Omega(m^3)$  lowerbound  $B_{pq}$ 2. main complications KDS

![](_page_30_Picture_2.jpeg)

![](_page_31_Figure_1.jpeg)

![](_page_32_Figure_1.jpeg)

![](_page_33_Figure_1.jpeg)

![](_page_34_Figure_1.jpeg)

![](_page_35_Figure_1.jpeg)

![](_page_36_Figure_1.jpeg)

![](_page_37_Figure_1.jpeg)

1.  $B_{pq}$  left to right  $\implies \Omega(m^2)$  events

![](_page_38_Figure_2.jpeg)

- 1.  $B_{pq}$  left to right  $\implies \Omega(m^2)$  events
- 2. Move  $B_{pq}$  left to right  $\Omega(m)$  times.

 $\implies \Omega(m^3)$  events

![](_page_39_Figure_4.jpeg)

- 1.  $B_{pq}$  left to right  $\implies \Omega(m^2)$  events
- 2. Move  $B_{pq}$  left to right  $\Omega(m)$  times.

```
\implies \Omega(m^3) events
```

![](_page_40_Figure_4.jpeg)

- 1.  $B_{pq}$  left to right  $\implies \Omega(m^2)$  events
- 2. Move  $B_{pq}$  left to right  $\Omega(m)$  times.

```
\implies \Omega(m^3) events
```

![](_page_41_Picture_4.jpeg)

- 1.  $B_{pq}$  left to right  $\implies \Omega(m^2)$  events
- 2. Move  $B_{pq}$  left to right  $\Omega(m)$  times.

```
\implies \Omega(m^3) events
```

![](_page_42_Figure_4.jpeg)

- 1.  $B_{pq}$  left to right  $\implies \Omega(m^2)$  events
- 2. Move  $B_{pq}$  left to right  $\Omega(m)$  times.

```
\implies \Omega(m^3) events
```

![](_page_43_Figure_4.jpeg)

- 1.  $B_{pq}$  left to right  $\implies \Omega(m^2)$  events
- 2. Move  $B_{pq}$  left to right  $\Omega(m)$  times.

```
\implies \Omega(m^3) events
```

![](_page_44_Picture_4.jpeg)

- 1.  $B_{pq}$  left to right  $\implies \Omega(m^2)$  events
- 2. Move  $B_{pq}$  left to right  $\Omega(m)$  times.

```
\implies \Omega(m^3) events
```

![](_page_45_Picture_4.jpeg)

- 1.  $B_{pq}$  left to right  $\implies \Omega(m^2)$  events
- 2. Move  $B_{pq}$  left to right  $\Omega(m)$  times.

```
\implies \Omega(m^3) events
```

![](_page_46_Picture_4.jpeg)

![](_page_47_Picture_1.jpeg)

![](_page_48_Picture_1.jpeg)

![](_page_49_Picture_2.jpeg)

![](_page_50_Picture_2.jpeg)

![](_page_51_Picture_2.jpeg)

![](_page_52_Picture_2.jpeg)

![](_page_53_Picture_2.jpeg)

1. Combined Events

![](_page_54_Picture_3.jpeg)

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![](_page_55_Picture_3.jpeg)

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![](_page_56_Picture_3.jpeg)

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![](_page_57_Picture_3.jpeg)

1. Combined Events

![](_page_58_Picture_3.jpeg)

1. Combined Events

![](_page_59_Picture_3.jpeg)

- 1. Combined Events
- 2. Implicit Bisector Certificates
- 3. Bisector Splits

![](_page_60_Figure_4.jpeg)

- 1. Combined Events
- 2. Implicit Bisector Certificates
- 3. Bisector Splits

![](_page_61_Figure_4.jpeg)

1. Combined Events

2. Implicit Bisector Certificates

3. Bisector Splits

**Thm.**  $\exists$  KDS O(m) size to maintain  $B_{pq}$  that handles events and updates in  $O(\log m)$  time each, and supports splits in  $O(\log^2 m)$  time.

![](_page_62_Picture_5.jpeg)

1. Combined Events

2. Implicit Bisector Certificates

3. Bisector Splits

![](_page_63_Picture_5.jpeg)

![](_page_63_Figure_6.jpeg)

1. Improve  $\tilde{O}(n^3m^3)$  bound on #events Can we get  $\Theta(nm^3 + n^2m)$ ?

**Thm.**  $\exists$  KDS O(n + m) size to maintain VD(S) that handles events in

![](_page_64_Picture_3.jpeg)

![](_page_64_Figure_4.jpeg)

Improve Õ(n<sup>3</sup>m<sup>3</sup>) bound on #events
 Can we get Θ(nm<sup>3</sup> + n<sup>2</sup>m) ?
events in O(polylog nm) time?

![](_page_65_Picture_3.jpeg)

![](_page_65_Picture_4.jpeg)

- 1. Improve  $\tilde{O}(n^3m^3)$  bound on #events Can we get  $\Theta(nm^3 + n^2m)$ ?
- 2. events in *O*(polylog *nm*) time?
- 3. more general movement?

![](_page_66_Picture_5.jpeg)

![](_page_66_Picture_6.jpeg)

- 1. Improve  $\tilde{O}(n^3m^3)$  bound on #events Can we get  $\Theta(nm^3 + n^2m)$ ?
- 2. events in *O*(polylog *nm*) time?
- 3. more general movement?
- 4. holes?

![](_page_67_Picture_6.jpeg)

![](_page_67_Picture_7.jpeg)